

EX1: Convert  $F = a \text{ xor } b$  into a BDD (ordering  $a < b$ )

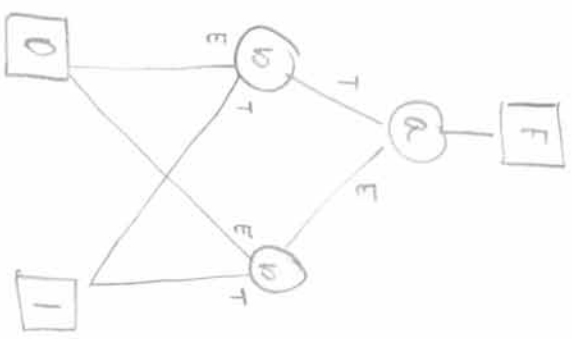
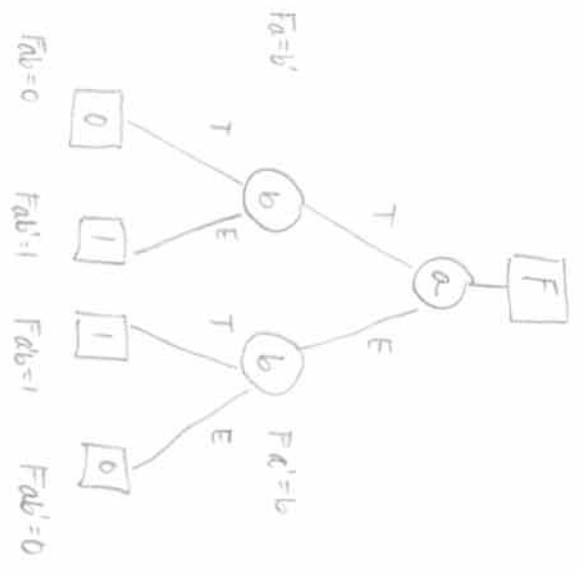
$$F = a'b + ab'$$

$$\swarrow \quad \searrow$$

$$F_{a=b'} \quad F_{a'=b}$$

$$\swarrow \quad \searrow$$

$$F_{ab} = 0 \quad F_{ab'} = 1 \quad F_{ab} = 1 \quad F_{ab'} = 0$$



you could use multiple terminal nodes in your drawing to make BDD easier to draw but internal structure will have 1 pair of terminal nodes

EX2: use ITE Algorithm 1 to convert  $F = ab + b'c'$  into BDD (variable ordering  $a \leq b \leq c$ )

According to ITE Operator table provided  $F + G = \text{ITE}(F, 1, G)$ , here  $F = ab$  and  $G = b'c'$

$$= \text{ITE}(ab, 1, b'c')$$

$$= \text{ITE}(a, \text{ITE}((ab)_a, (1)_a), \text{ITE}(b'c')_a, (1)_{a'})$$

$$= \text{ITE}(a, \text{ITE}(0, 1, b'c'), \text{ITE}(b, 1, b'c')) \quad // \text{ITE}(0, 1, b'c') = b'c'$$

$$= \text{ITE}(a, b'c', \text{ITE}(b, 1, b'c')) \quad // \text{expand } \text{ITE}(b, 1, b'c') \text{ w.r.t. } b$$

$$= \text{ITE}(a, b'c', \text{ITE}(b, \text{ITE}(b)_b, (1)_b), \text{ITE}(b)_b, (1)_{b'})$$

$$= \text{ITE}(a, b'c', \text{ITE}(b, \text{ITE}(1, 1, 0), \text{ITE}(0, 1, c')))$$

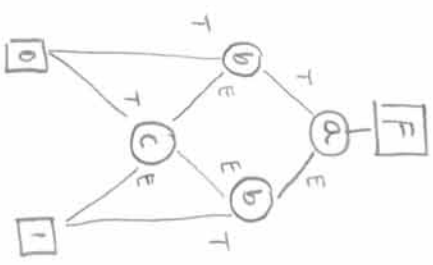
$$= \text{ITE}(a, b'c', \text{ITE}(b, 1, c')) \quad // b'c' = \text{ITE}(b', c', 0) \text{ according to table provided.}$$

$$= \text{ITE}(a, \text{ITE}(b', c', 0), \text{ITE}(b, 1, c')) \quad // \text{cannot have a node labeled } b' \text{ so we expand w.r.t } b$$

$$= \text{ITE}(a, \text{ITE}(b, \text{ITE}(b)_b, (c')_b), \text{ITE}(b)_b, (c')_{b'})$$

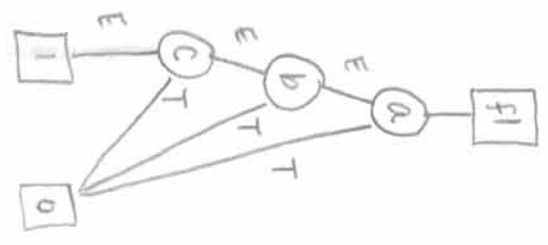
$$= \text{ITE}(a, \text{ITE}(b, \text{ITE}(0, c', 0), \text{ITE}(1, c', 0)), \text{ITE}(b, 1, c'))$$

$$= \text{ITE}(a, \text{ITE}(b, 0, c'), \text{ITE}(b, 1, c'))$$

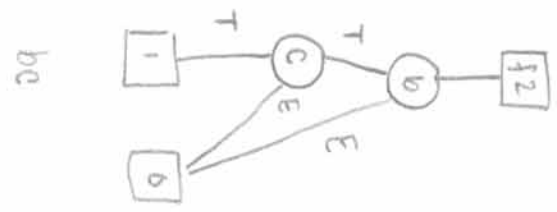


ex3: Iterative Approach to building BDDs.

$$F = a'b'c' + bc \quad (a \leq b \leq c)$$

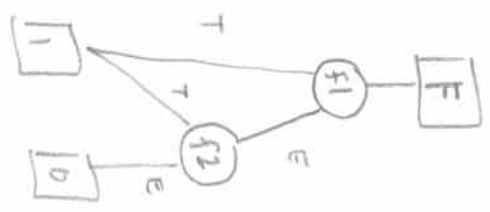


a'b'c'

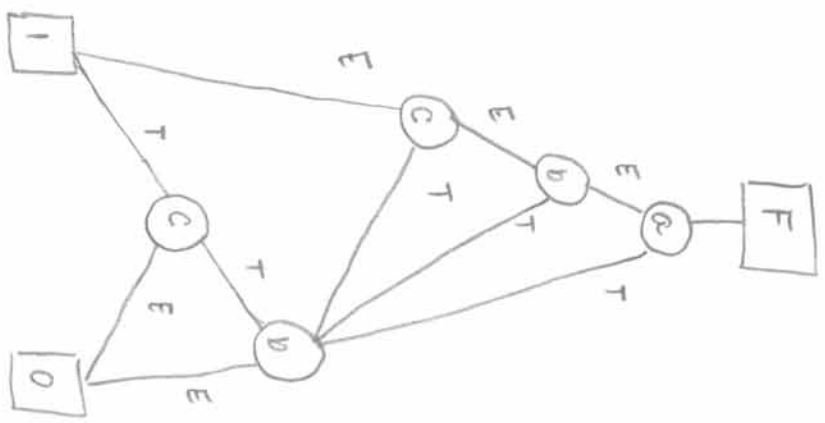


bc

How do I combine them?  
 $F = f1 + f2$



expand f1 & f2



problem! lost variable  
 ordering

no longer  
 reduced.

try again in ex4 with  
 recursive approach.

Ex 4: ITE Recursive Approach to BDD's

$$F = a'b'c' + bc \quad (a \leq b \leq c)$$

$$\begin{aligned}
 &= \text{ITE}(a'b'c', 1, bc) \\
 &= \text{ITE}(a, \text{ITE}(a'b'c', a, (1)_a, (bc)_a), \text{ITE}(a'b'c', (1)_a, (bc)_a)) \\
 &= \text{ITE}(a, \text{ITE}(0, 1, bc), \text{ITE}(b'c', 1, bc)) \\
 &= \text{ITE}(a, bc, \text{ITE}(b, \text{ITE}(b'c', (1)_b, (bc)_b), \text{ITE}(b'c', (1)_b, (bc)_b))) \\
 &= \text{ITE}(a, bc, \text{ITE}(b, \text{ITE}(0, 1, c), \text{ITE}(c', 1, 0))) \\
 &= \text{ITE}(a, bc, \text{ITE}(b, c, \text{ITE}(c, \text{ITE}(0, 1, 0), \text{ITE}(1, 1, 0)))) \\
 &= \text{ITE}(a, bc, \text{ITE}(b, c, \text{ITE}(c, 0, 1))) \\
 &= \text{ITE}(a, bc, \text{ITE}(b, c, c')) \\
 &= \text{ITE}(a, \text{ITE}(b, c, 0), \text{ITE}(b, c, c'))
 \end{aligned}$$

✓ ordered  
✓ reduced.

